# Challenges of Serverless: can languages help?

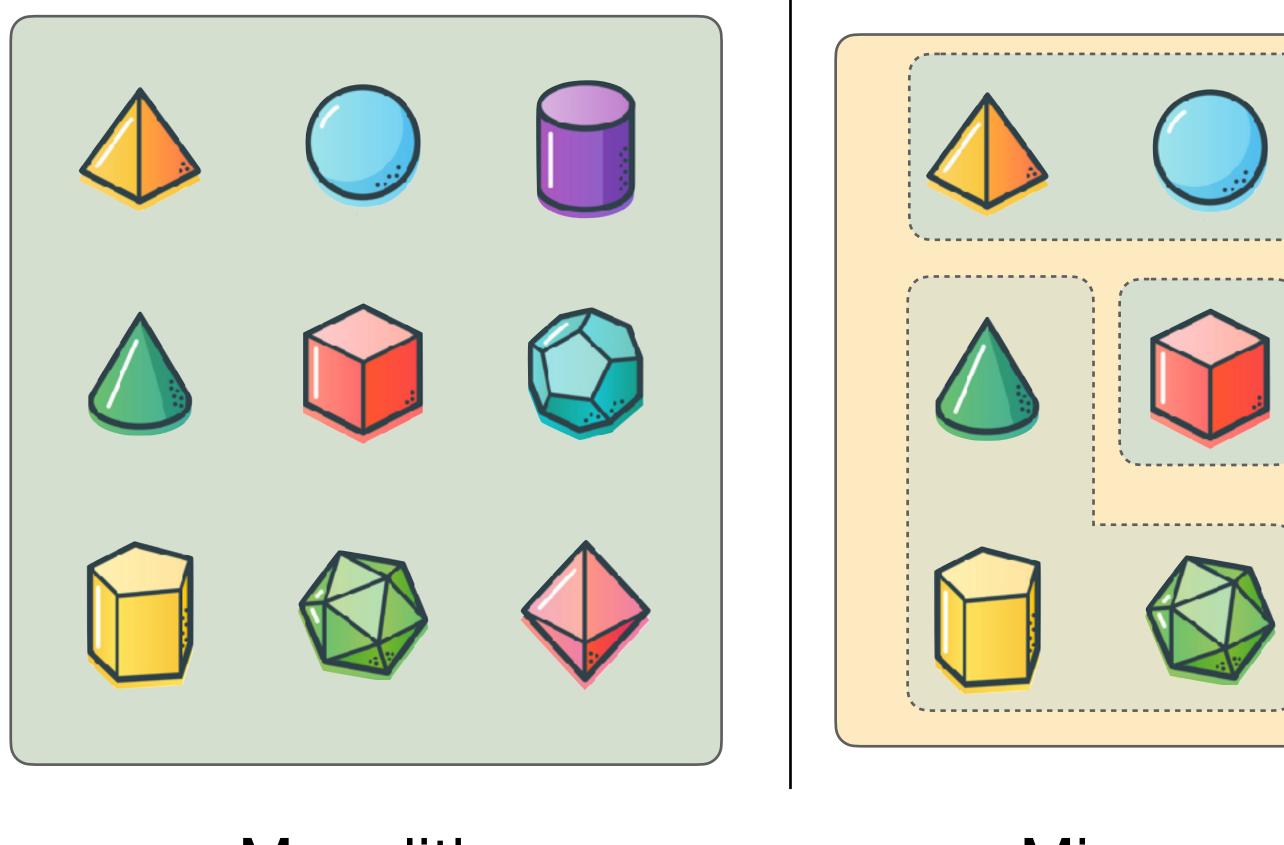
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### Serverless (and Microservices)

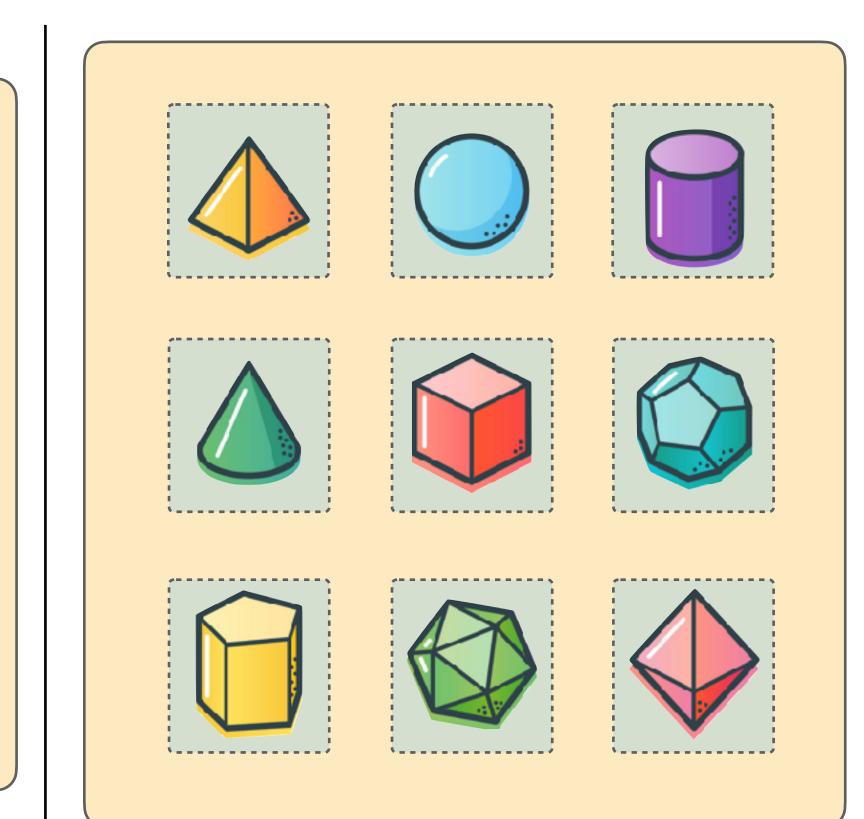
### provisioned, pay-per-deployment



### Monolith



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### on-demand, pay-per-execution

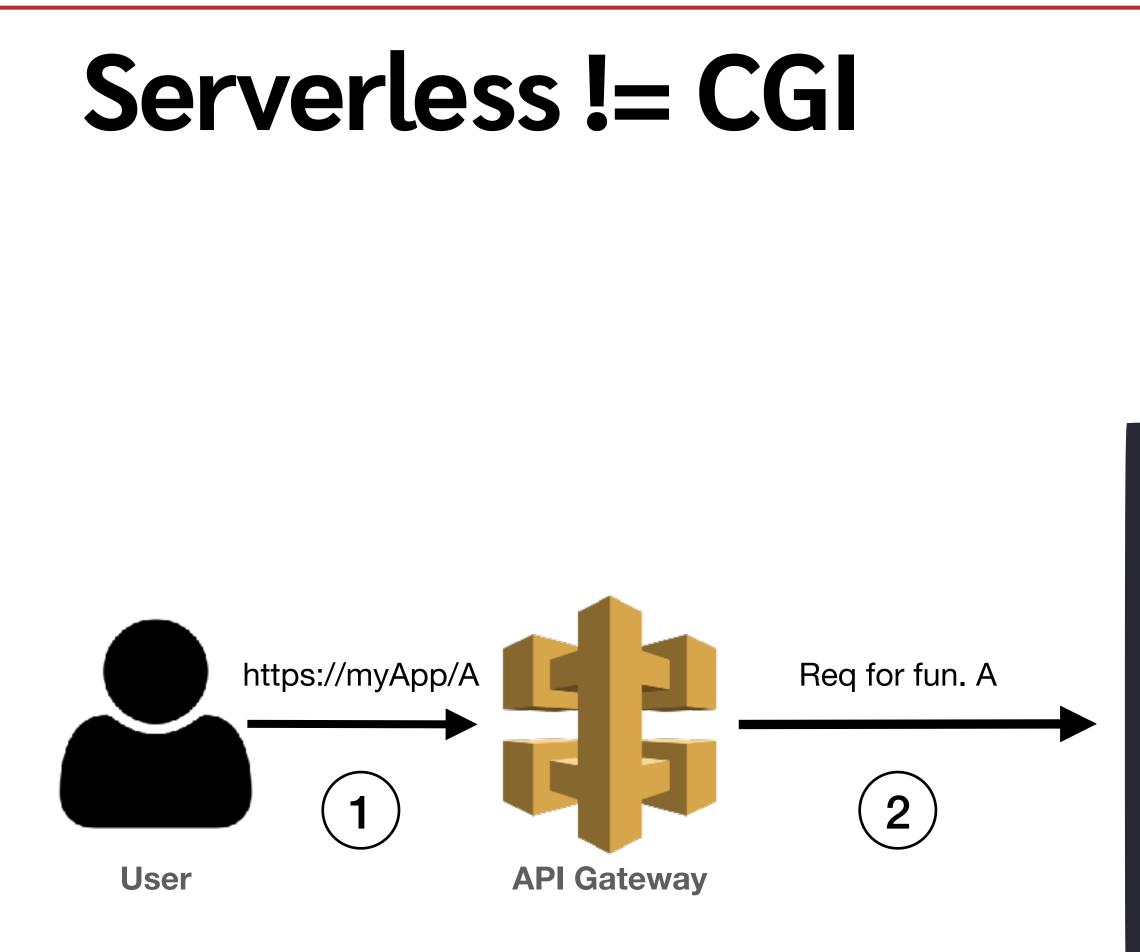
### Microservices

### Serverless



Runtime Environment



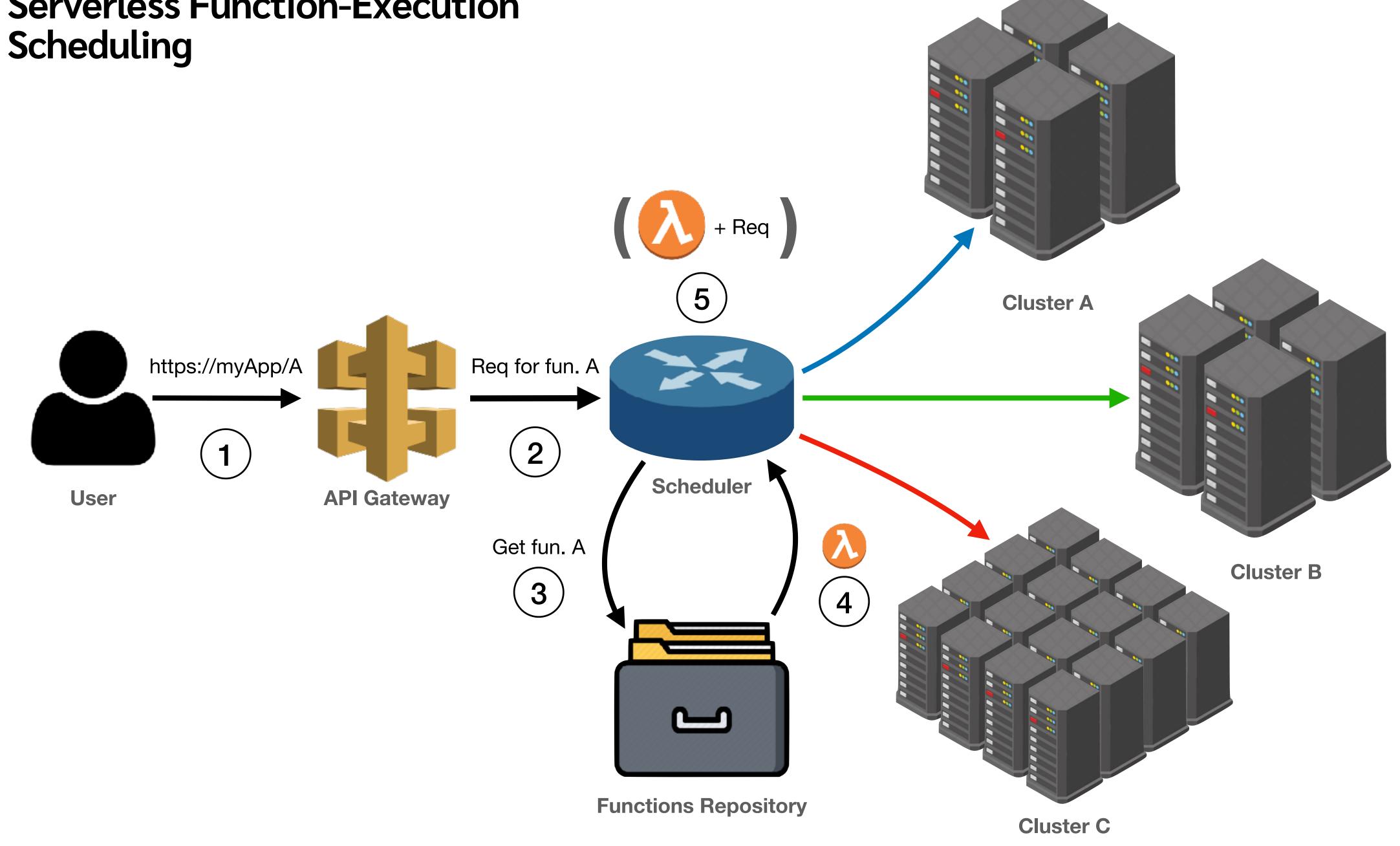


```
# A serverless cgi-bin!
# https://www.hawksworx.com/cgi-bin/hello/friend
[[redirects]]
from = "/cgi-bin/hello/:name"
to = "/.netlify/functions/hello?name=:name"
status = 200
```





## **Serverless Function-Execution**

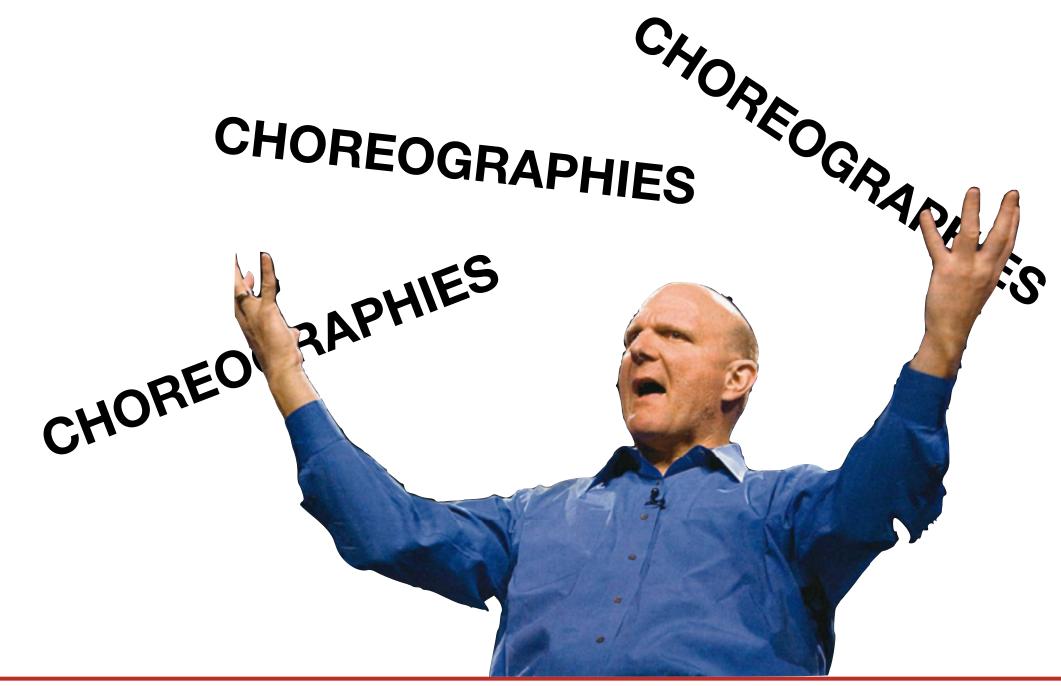


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Few programmers can correctly write down their program semantics in a sequential language while also accounting for **parallel interleaving**, **message reordering**, **partial failures**, and **dynamic autoscaling** deployment.

Availability protocols are frequently interleaved into program logic in ways that make them tricky to test and evolve.

From "New Directions in Cloud Programming"





State management/security: cloud applications often need to share or exchange short-lived/ephemeral state among its components, e.g., application-wide caches, indexes, lookup tables, intermediate results. In serverless, this is usually solved via object storage and key-value stores, but the logic of the distributed system becomes fragmented, even when we neglect to consider access control, availability, consistency, etc.

One possibility is to have stateful cloud functions (e.g., wrapping caches)[SKC]

From "What Serverless computing is and should become"

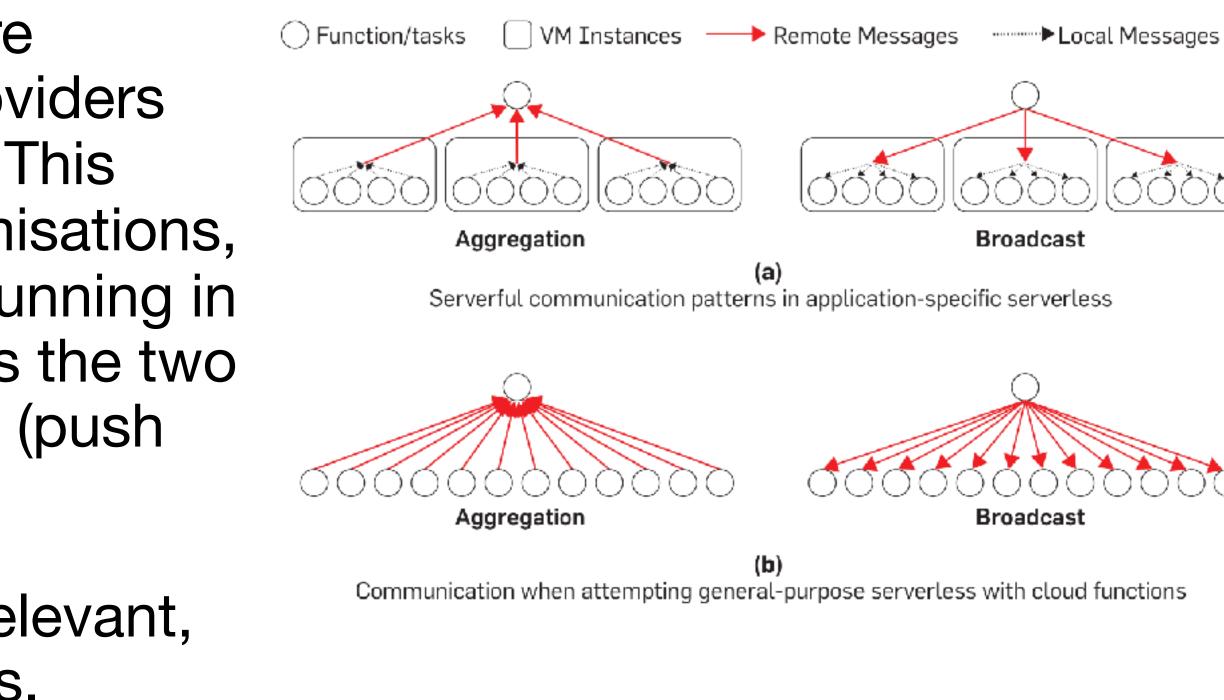


Scheduling: users cede control over where functions run and, as a consequence, providers cannot support direct communication. This prevents users to express serverful optimisations, e.g., passing state between functions running in the same machine, which instead requires the two functions to run through at least two trips (push and pull) from some shared storage.

The shortcomings of this become more relevant, e.g., in the case of scatter/gather patterns.

A solution here is to let languages express this information (from developers) and provide it to systems like APP so they can figure out the most efficient (as in, cheapest, fastest, greenest, etc.) way to achieve that.

From "What Serverless computing is and should become"



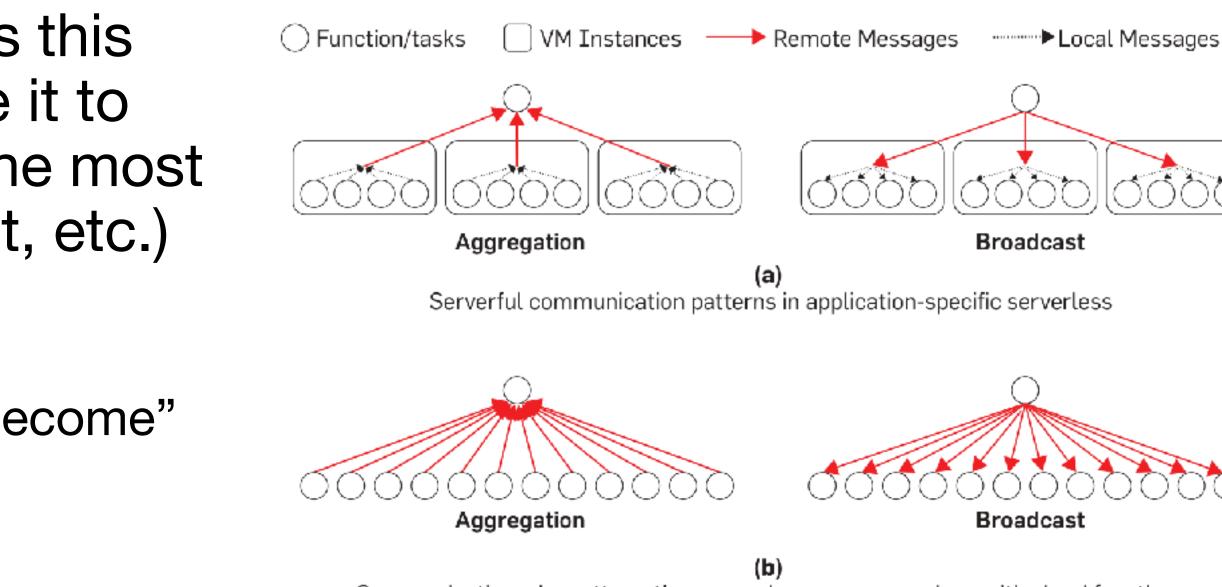






A solution here is to let languages express this information (from developers) and provide it to systems like APP so they can figure out the most efficient (as in, cheapest, fastest, greenest, etc.) way to achieve that.

From "What Serverless computing is and should become"



Communication when attempting general-purpose serverless with cloud functions

...an environment where programmers can specify multi-objective performance targets for execution, e.g., tradeoffs between billing costs, latency, and availability.

From "New Directions in Cloud Programming"









```
Simple COVID-19 Tracker App in Pythonic HydroLogic
   class Person: (pid: int, country: string,
2
           contacts: Set(&Person), covid: bool, vaccinated: bool
           key=pid, partition=country)
   table people: Person
4
5
   var vaccine_count: int
   on add_person(pid: int):
     people.merge(Person(pid)) # monotonic mutation
8
     return OK
9
   on add_contact(p: Person, p1: Person):
     p.contacts.merge(p1) # monotonic mutation
     p1.contacts.merge(p) # monotonic mutation
     return OK
16
   query transitive(p: Person, p1: Person): # monotonic query
       {(p, p1) for p in people for p1 in p.contacts}
       {(p, p2) for (p, p1) in transitive for p2 in p1.contacts}
8
   on trace(p: Person):
     return (p2 for (p, p2) in transitive(p, _)
```

### Hydrologic, from "New Directions in Cloud Programming"

```
on diagnosed(pid: int):
     people[pid].covid.merge(true) # monotonic mutation
     send alert(p: Person) {p for p in trace(pid)}
   from covid_xmission_model import covid_predict
   on likelihood(pid: int):
     return covid_predict(people[pid])
    on vaccinate(pid: int, consistency={serializable;
                   vaccine_count >= 0; people.has_key(pid)}):
     people[pid].vaccinated.merge(True) # monotonic mutation
     vaccine_count := vaccine_count - 1 # NON-monotonic mutation
      return OK
   availability:
Β7
        default: { domain = AZ, failures = 2 }
38
        likelihood: { domain = AZ, failures = 1 }
   target:
41
        default: { latency = 100ms, cost = 0.01units }
42
        likelihood: { processor = GPU, cost = 0.1units }
```





```
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```

Hydrologic considers an event loop that runs on successive snapshots of the overall program state, including new inbound messages to be handled.

Each iteration of the loop uses the developer's program specification to compute new results from the snapshot and atomically updates the state of the program and the end of the loop.

This is a kind of "single-node" model that leaves to other part of the specification (bottom) the task to define their deployment-related information.

E.g., they can indicate a certain number of endpoints for a specific function and define load-balancing and replication policies to meet a given SLA on response rate.

Hydrologic, from "New Directions in Cloud Programming"



- one? Could the language help in understanding/capturing the characteristics of serverless and serverful architectures?
- Can we use choreographies to define a server\* architecture, e.g., specifying/constraining parts of the interactions with a specific for the other ones (e.g., sets of messages, queues, etc)
  - proxy services?

• Can we use Jolie to deploy a microservices architecture as a serverless

semantics, but leave the compiler/deployer decide the "best" strategies

• For the constrained parts, can we replicate (as in "preserve") the platform-specific semantics in either setting, e.g., by synthesising

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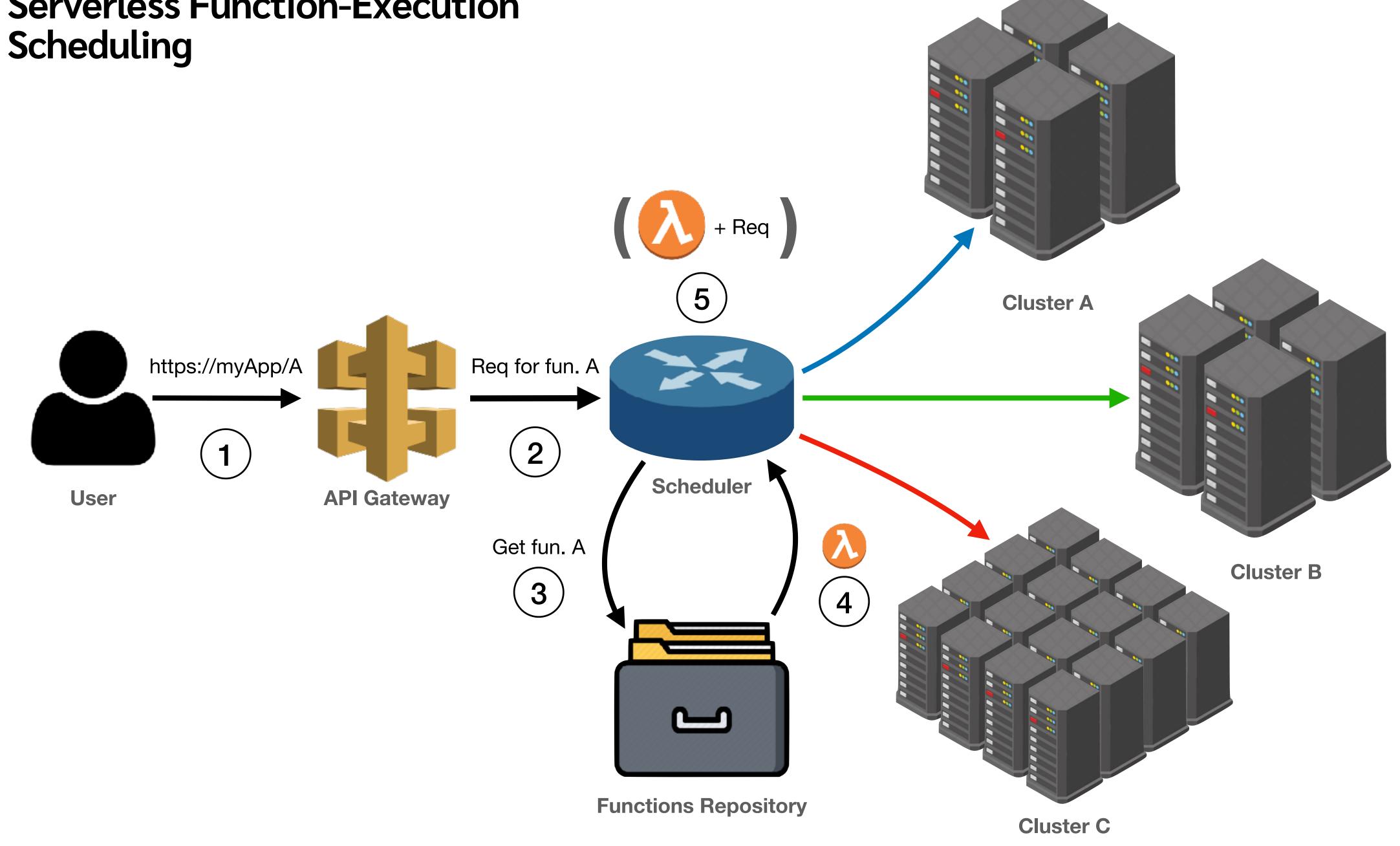
## Allocation Priority Policies for Serverless Function-execution Scheduling Optimisation

Giuseppe de Palma<sup>1</sup>, <u>Saverio Giallorenzo<sup>1,2</sup></u>, Jacopo Mauro<sup>3</sup> and Gianluigi Zavattaro<sup>1,2</sup>

<sup>1</sup>Università di Bologna (IT) <sup>2</sup>INRIA (FR) <sup>3</sup>University of Southern Denmark (DK)

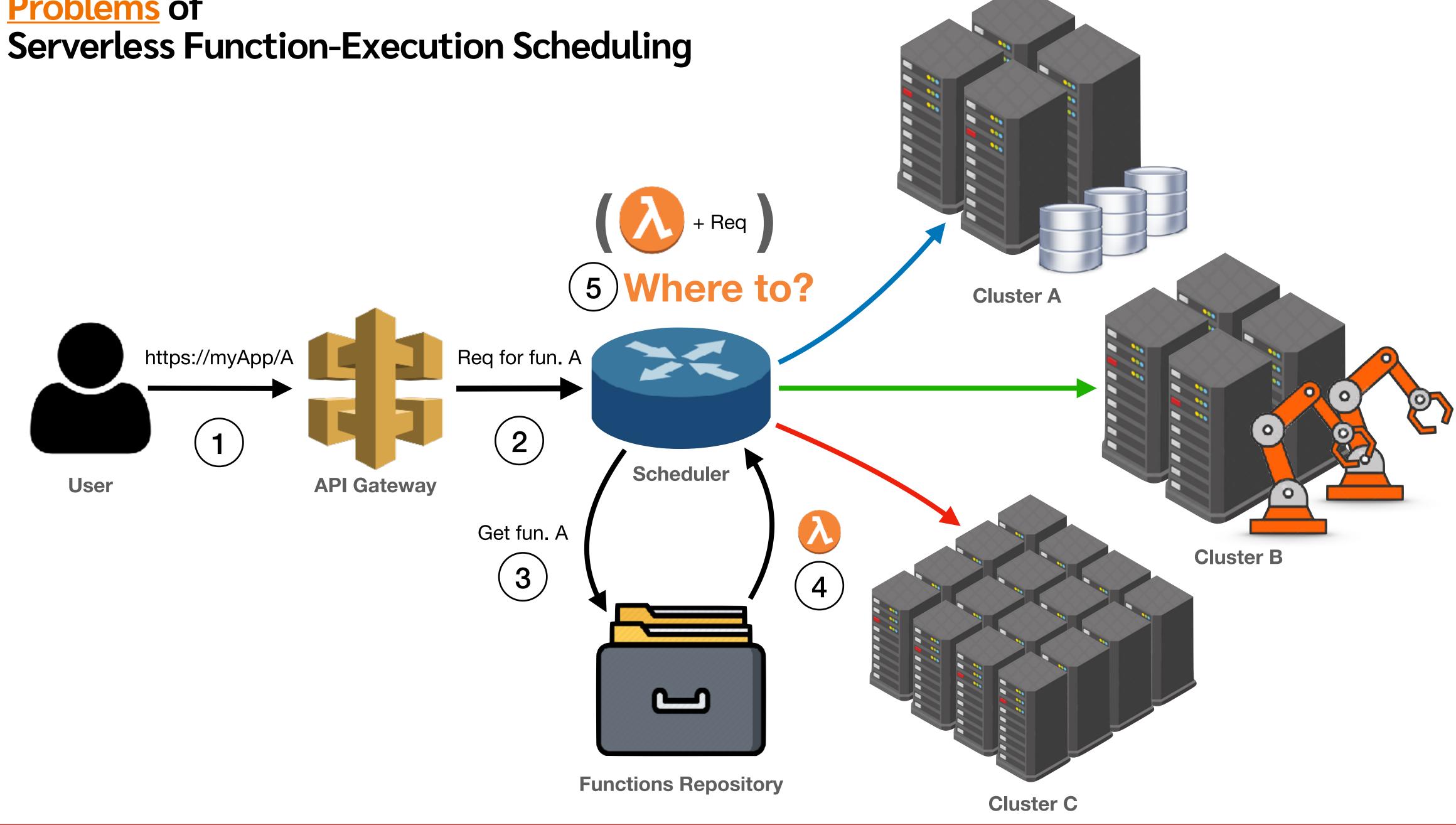


## **Serverless Function-Execution**





## **Problems** of

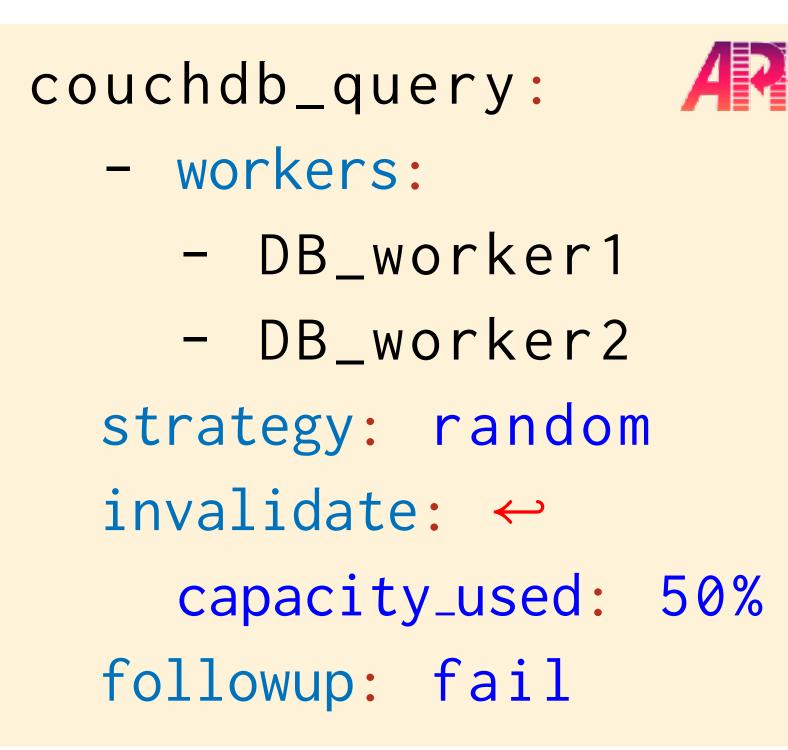


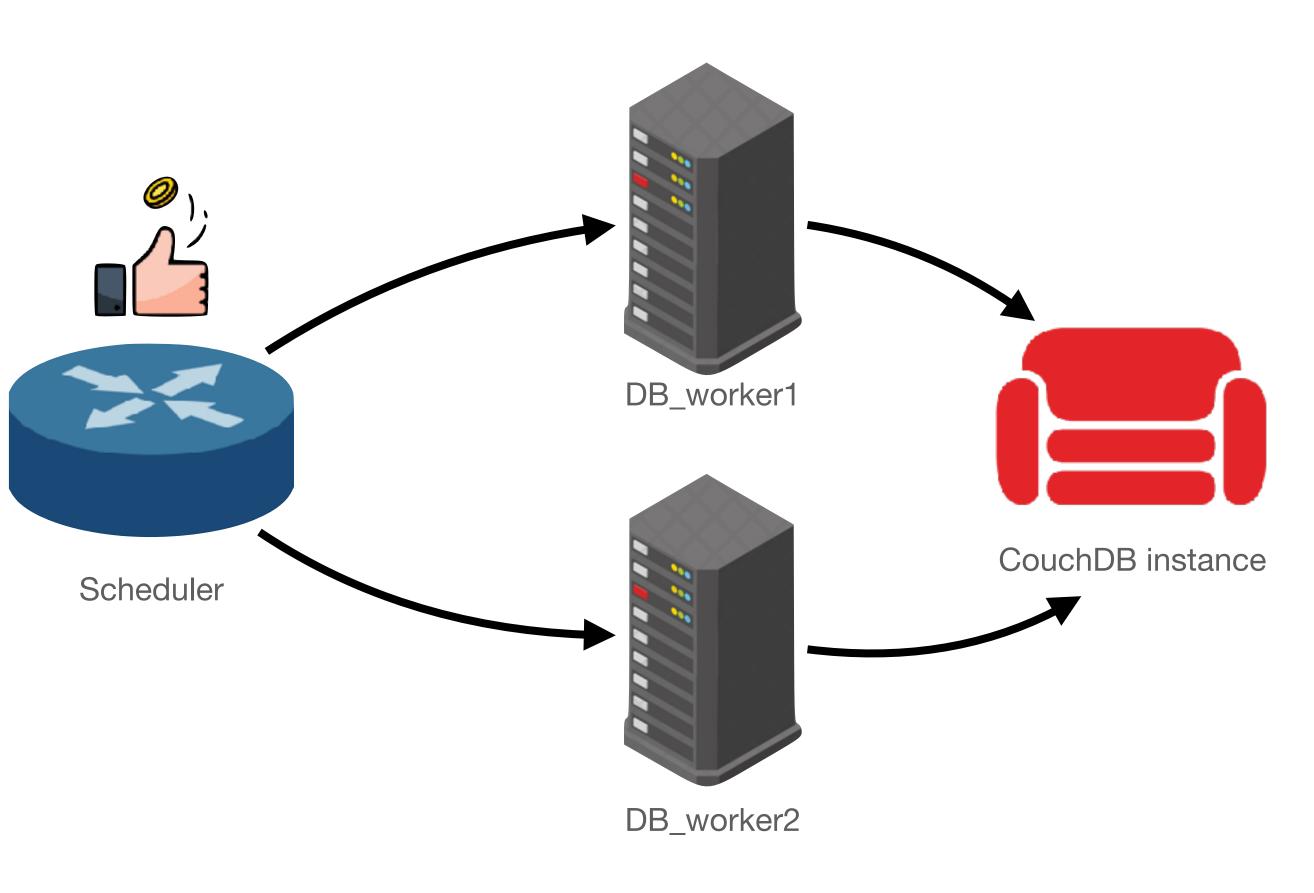
@ Artificial Intelligence, Cybersecurity, and Programming Languages group





### The APP Language • First Example



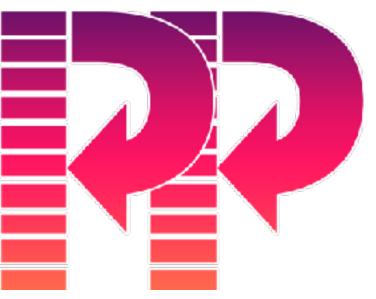






## The APP Language • Syntax

$policy\_tag$	e	$Identifiers \cup \{default\}$
app	::=	$\overline{tag}$
tag	::=	policy_tag : - block followup?
block		<pre>workers [ "*"   - worker_label (strategy [ random   platform (invalidate [ capacity_used : r</pre>
followup	::=	<pre>followup : [ default   fail ]</pre>



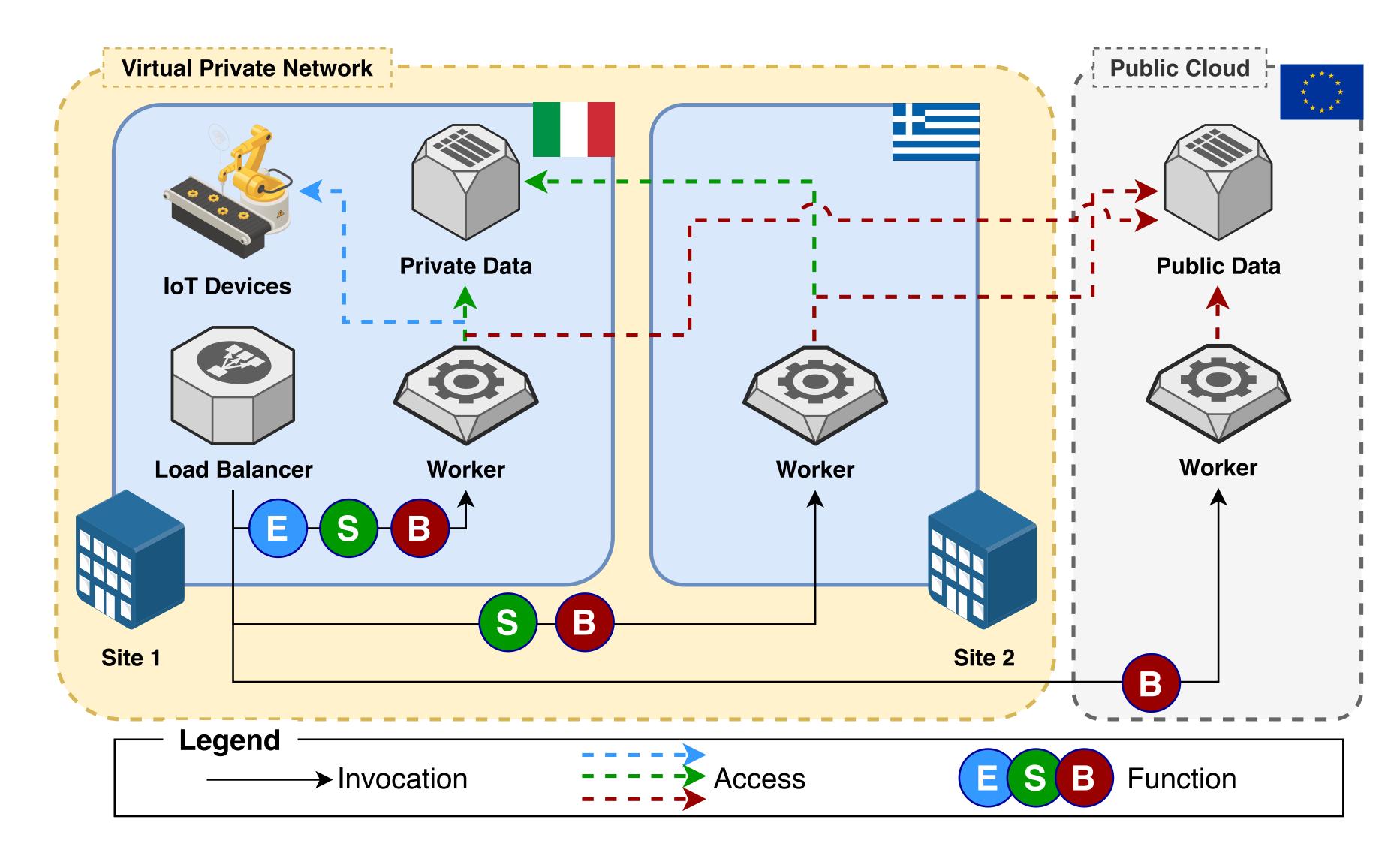
### $worker\_label \in Identifiers$ $n \in \mathbb{N}$





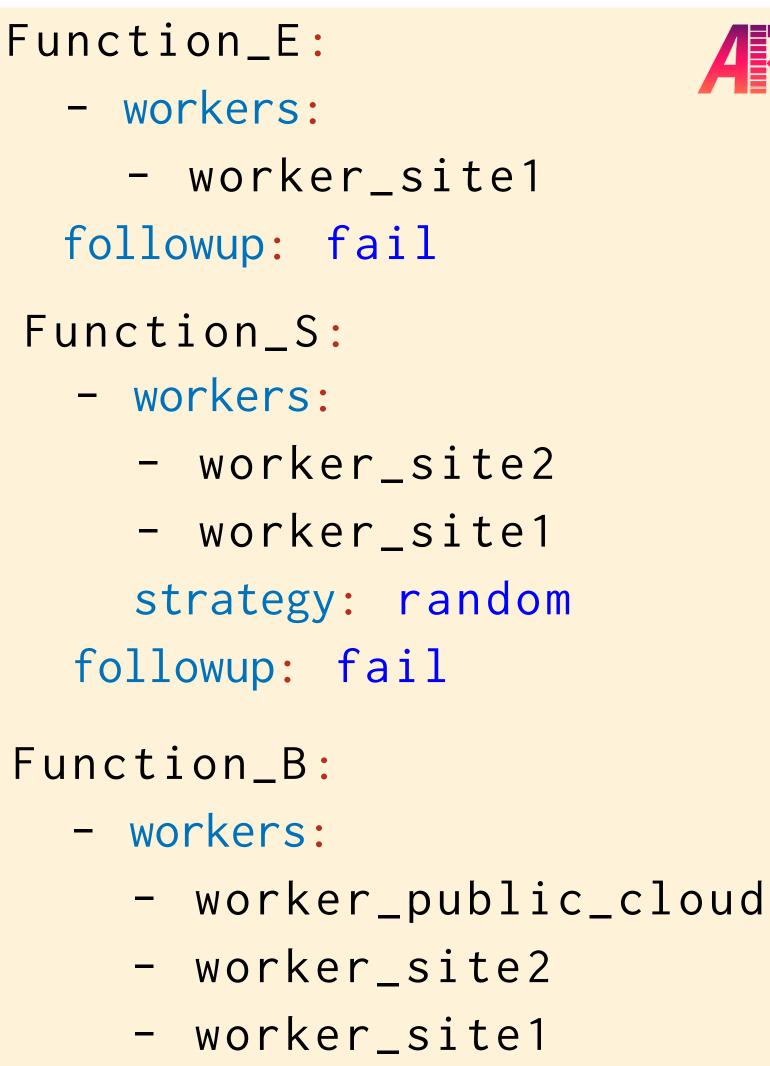


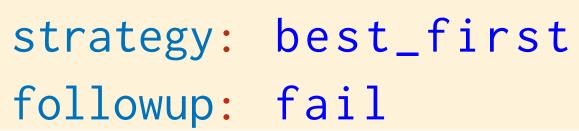
### Use case



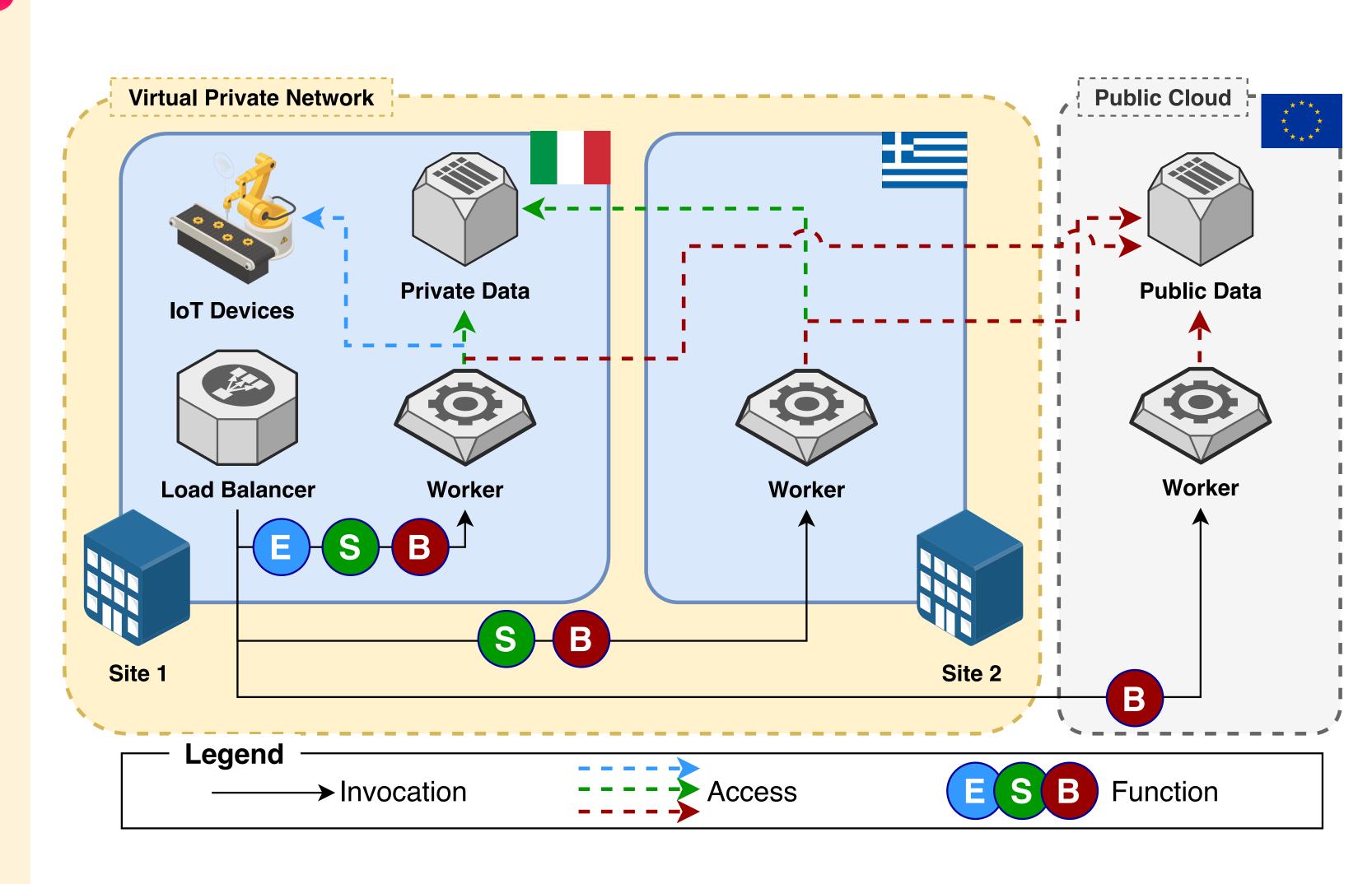
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### Use case - the APP deployment











## Use case - empirical results

	Site 1	Site 2	Public Cloud	Average (ms)	95% Average (ms)
E	1000	0	0	1096.53	1019.03
S	466	534	0	149.18	90.86
B	0	90	910	105.18	64.62

Table 1. 1000 invocation for each function in the APP-based OpenWhisk deployment.

		Site 1	Site 2	Public Cloud	Average (ms)	95% Average (ms)
OW1	E	1000	0	0	1159.90	1025.52
OW2	S	19	981	0	385.30	302.08
OW3	B	185	815	0	265.69	215.793

Table 2. 1000 invocations for each function in the vanilla OpenWhisk deployment.



### Future Work

- Automatic configuration of priority policies (ML, heuristics, etc.);
- Extend our prototype to support pools of workers;
- Test the expressiveness of APP by capturing and implementing the policies presented other papers on Serverless scheduling;
- Extend APP to describe (and not just use) scheduling algorithms and support the creation of user-defined libraries;
- Formalise the semantics of APP, useful for both a rigorous specification and to automatically reason on the properties of APP-defined deployments.





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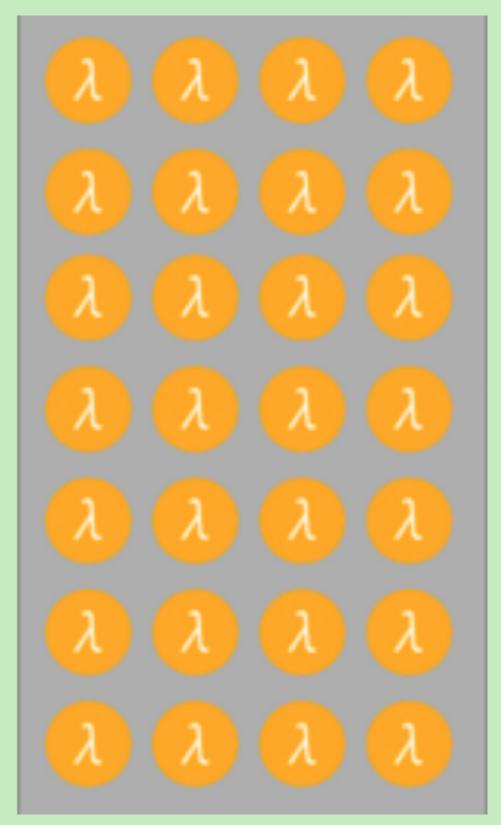


### Serverless (and Microservices)

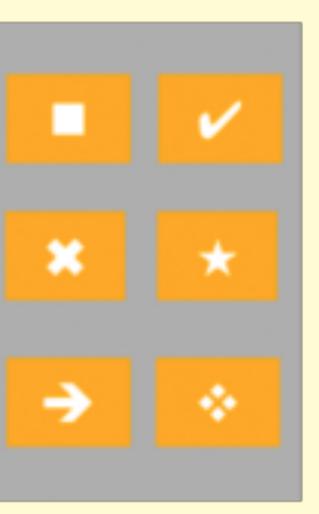
provisioned, pay-per-deployment  $\star$ \* О

Monolith









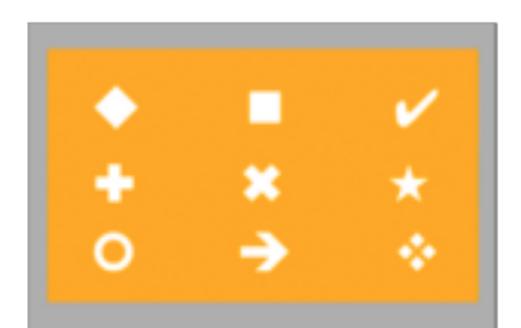
**Microservices** 







### Serverless (and Microservices)

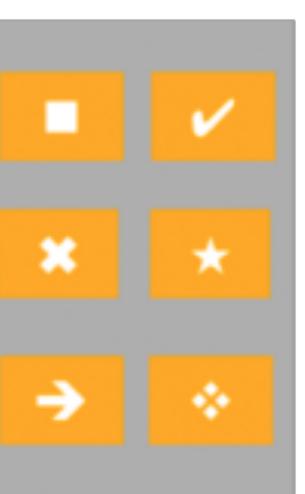


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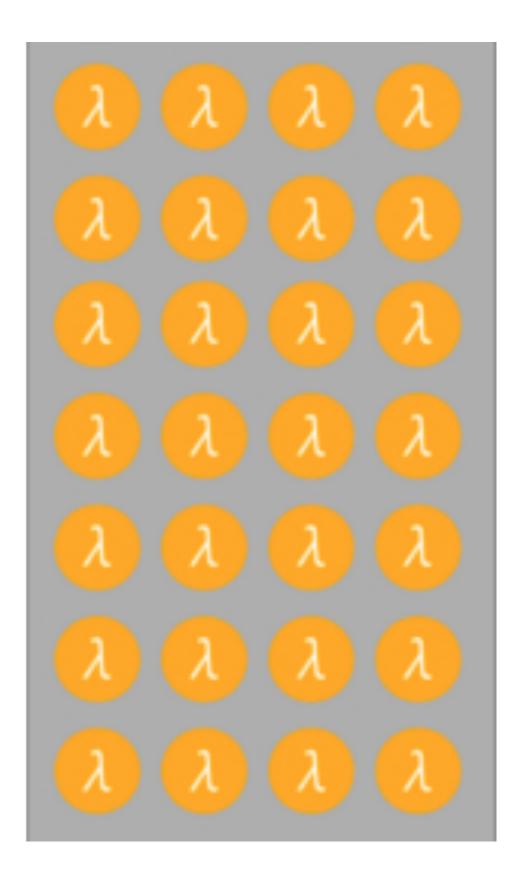
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### **Microservices**



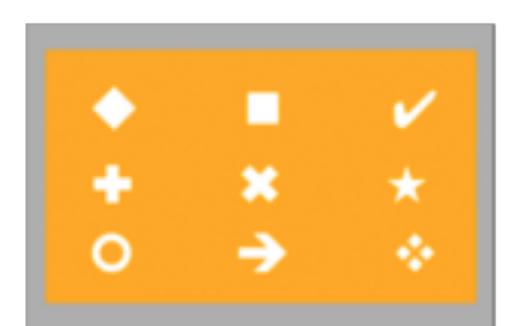








### **Serverless (and Microservices) • Readiness**

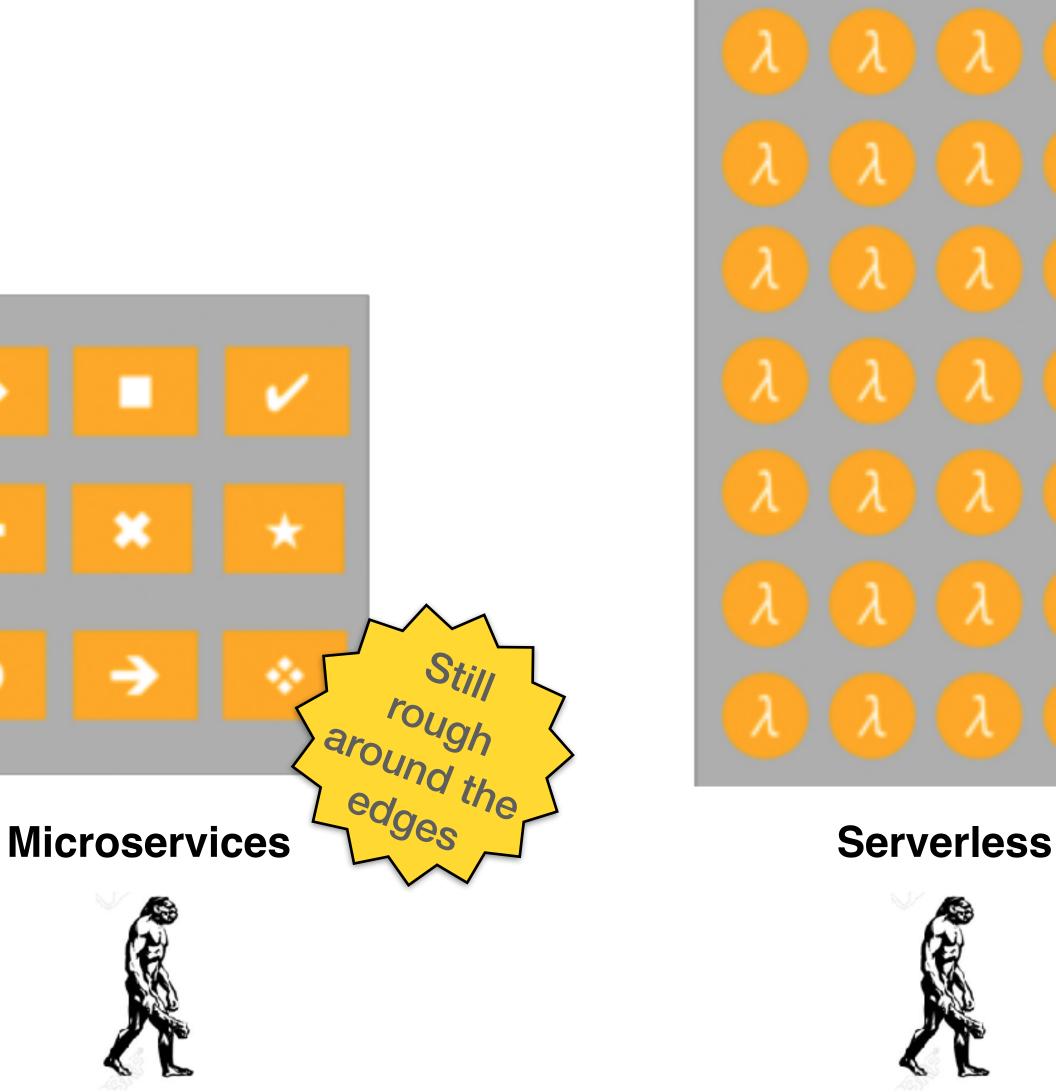


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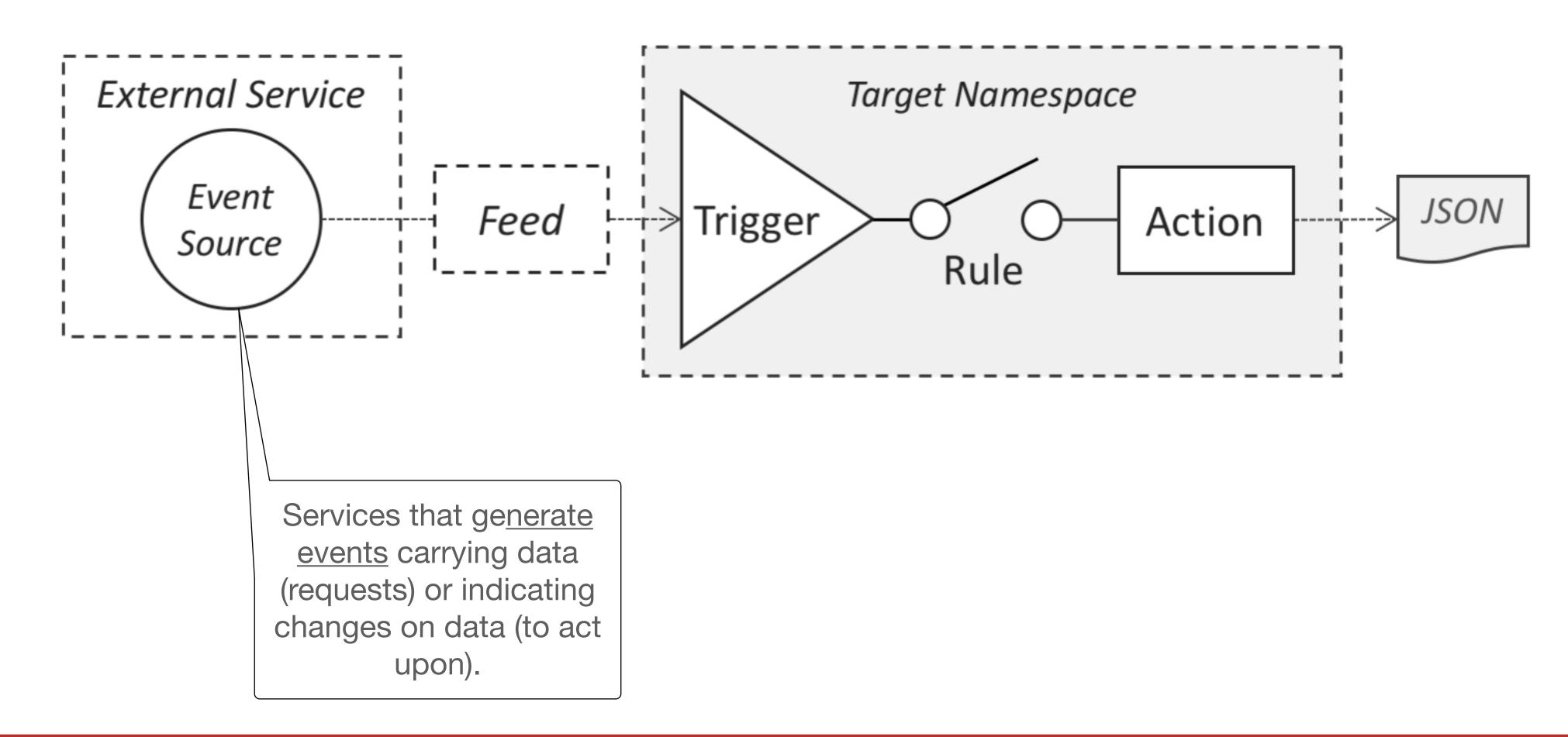
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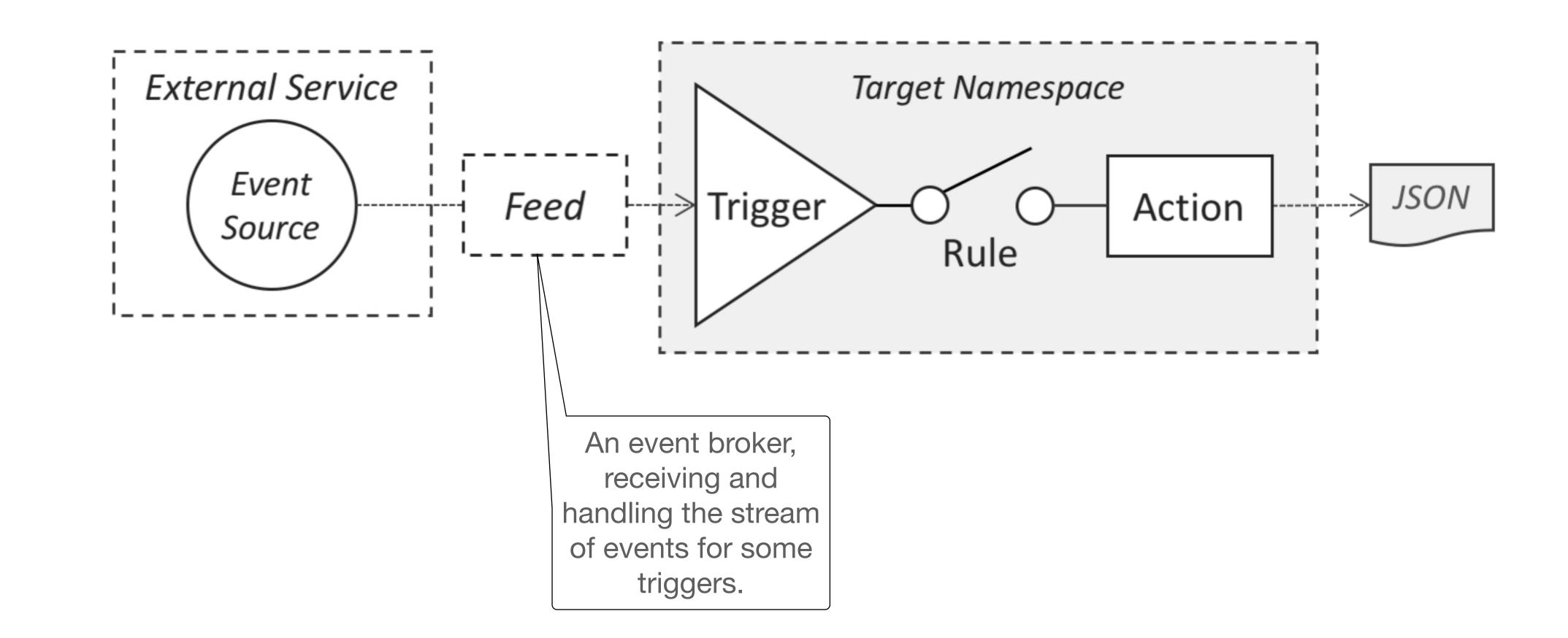




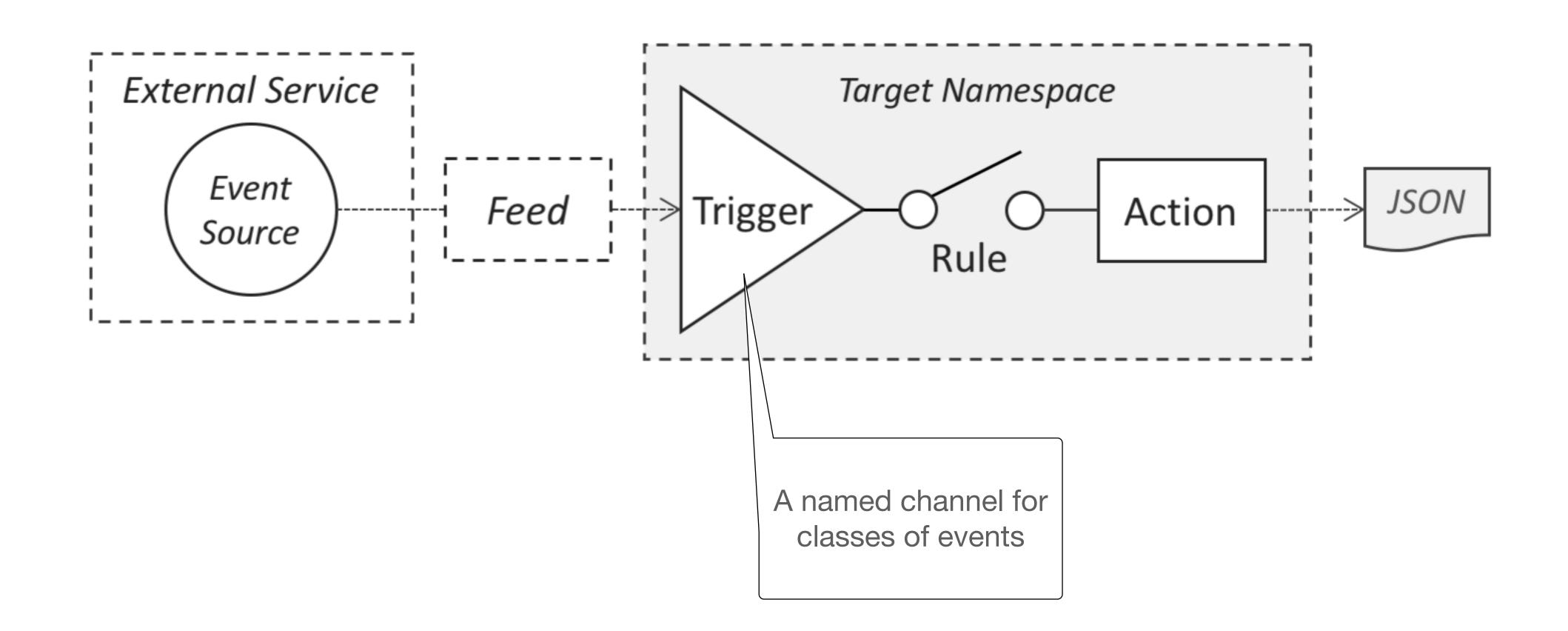




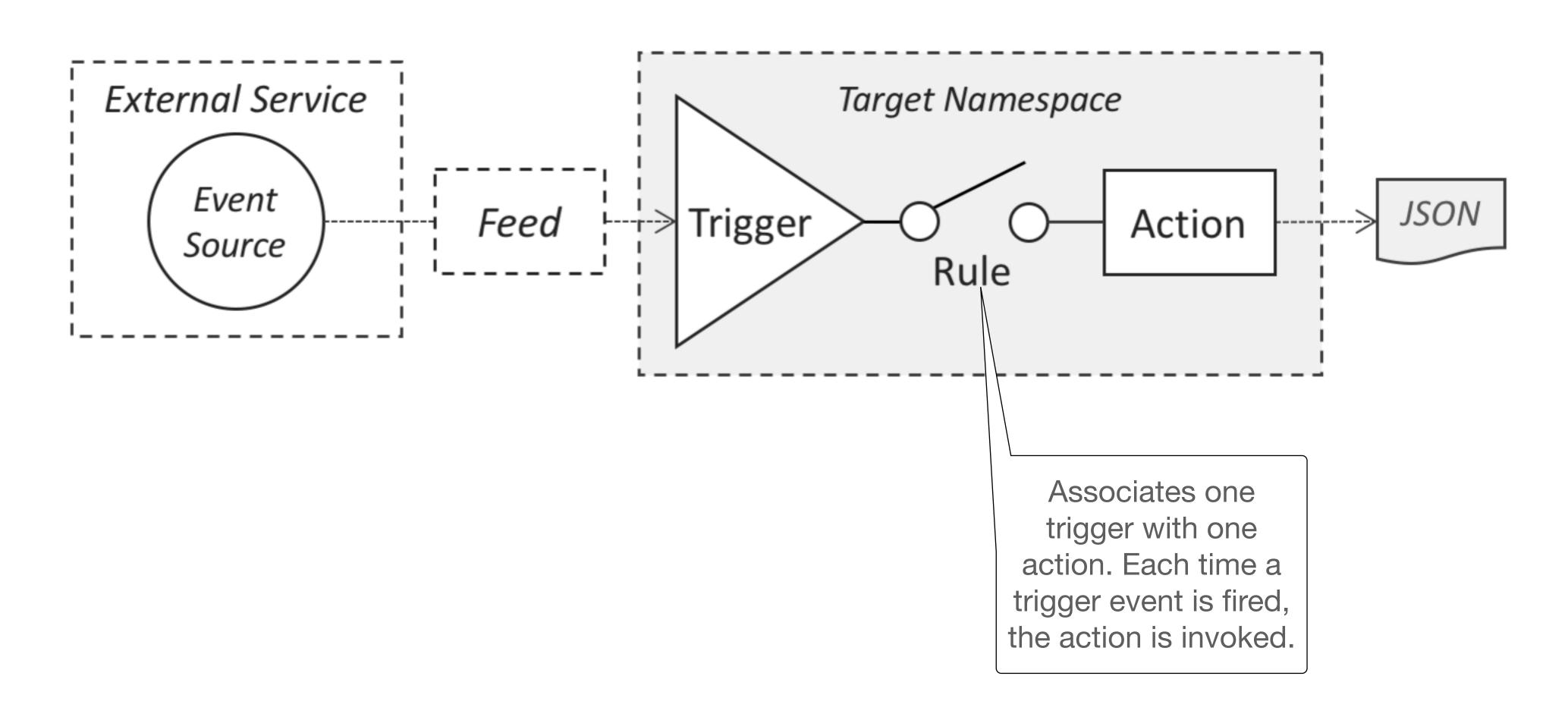






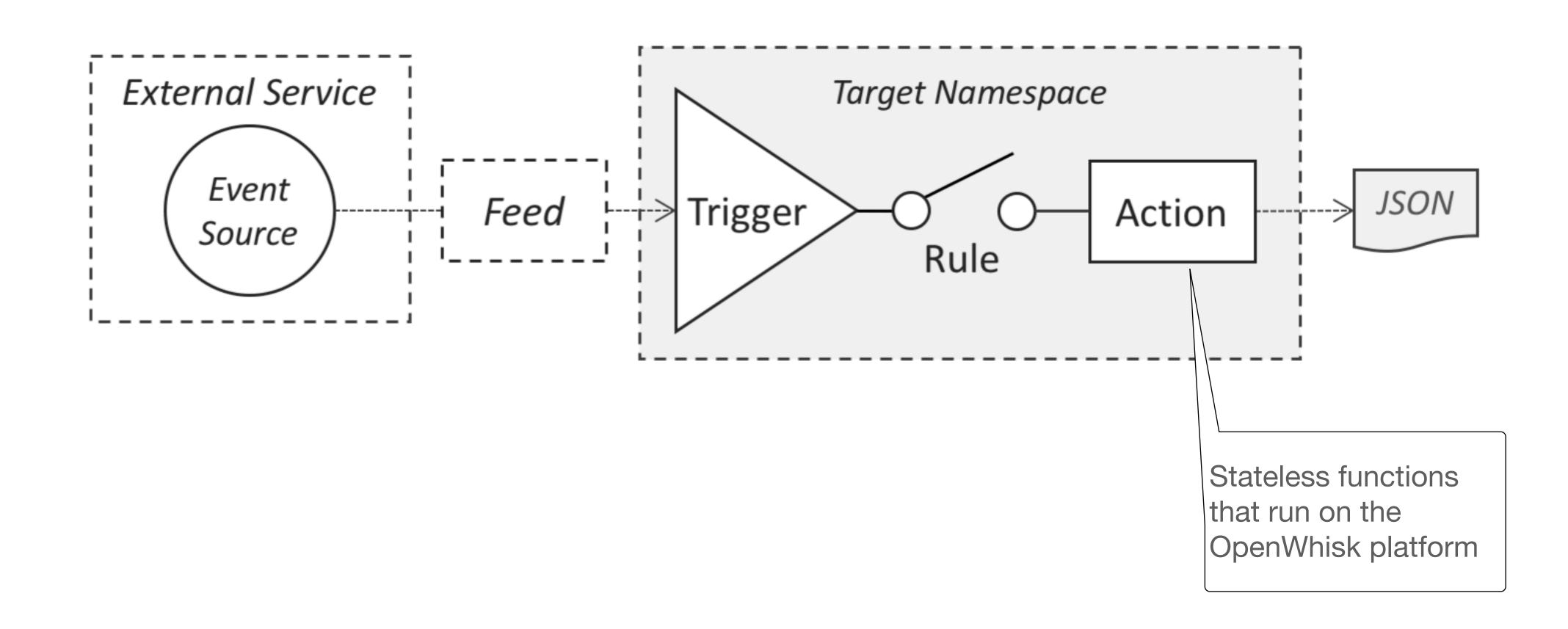














## Apache OpenWhisk

