# Serverless Computing in OLAS: Declarative Scheduling and Edge Cloud Platforms

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# A Gentle Introduction to Serverless





















### The APP Language • an example



De Palma, G., Giallorenzo, S., Mauro, J., & Zavattaro, G. (2020). Allocation Priority Policies for Serverless Function-Execution Scheduling Optimisation. In E. Kafeza, B. Benatallah, F. Martinelli, H. Hacid, A. Bouguettaya, & H. Motahari (Eds.), Service-Oriented Computing - 18th International Conference, ICSOC 2020, Dubai, United Arab Emirates, December 14-17, 2020, Proceedings (pp. 416–430). Springer.





## The APP Language • Case Study



- workers:
  - worker\_site1
- followup: fail
- Function\_S:
  - workers:
    - worker\_site2
    - worker\_site1
    - strategy: random
  - followup: fail
- Function\_B:
  - workers:
    - worker\_public\_cloud
    - worker\_site2
    - worker\_site1
  - strategy: best\_first
    followup: fail







## The APP Language • Case Study, results

	Site 1	Site 2	Public Cloud	Average (ms)	95% Average (ms)
E	1000	0	0	1096.53	1019.03
S	466	534	0	149.18	90.86
B	0	90	910	105.18	64.62

Table 1. 1000 invocation for each function in the APP-based OpenWhisk deployment.

		Site 1	Site 2	Public Cloud	Average (ms)	95% Average (ms)
OW1	E	1000	0	0	1159.90	1025.52
OW2	S	19	981	0	385.30	302.08
OW3	B	185	815	0	265.69	215.793

Table 2. 1000 invocations for each function in the vanilla OpenWhisk deployment.



, and for each service invocation call  $h(\overline{E})$ , we assume that h is a correct service verless invocations, which possibly execute on a different worker of the caller. niSL by means of three examples. As a first example, consider the code in Listing Perverless Scheduling Policies based on ( of a function that selects a functionality based on the characteristic of the involver. verless Scheduling Policies based on Cost if( icPromiumlcor // name: lambda1.miniSL // tag: premUser 1V/7( isPremiumUser, par ) => { if( isPremiumUser ) { 4 call PremiumService( par ) } else { call BasicService( par ) g 1! Function with a conditional statement of ke eithera PremiumService or a Başieservice depending on whether it has been ium user or  $main(u, v, P, B) = if_2(u, P, B)$ u = 1] a population ember (when  $if_2(u, P, B) = P$ be forwarde  $if_2(u, P, B) = B$  $\begin{bmatrix} u = 0 \end{bmatrix}$ be forwarde  $I_2(u, P, B) = B$  [u = 0] storting paper threes ample in the second set that checks the username and returns we want to reduce the latency of this function the best follow to the the present of the second set this case, it is difficult to predict the latency of the invocation of either the second set of the set of the set of the branch that will be selected is not known at and the latency of the invocation of either the second set of the set of the set of the branch the branch the selected is not known at and the latency of the invocation of either the set of the set of the set of the set of the branch the branch the selected is not known at and the set of the selected is not known at the set of the set of the set of the branch the branch the branch the selected is not known at and the set of the selected is not known at the set of the set of the set of the set of the branch the set of the set o



Inference of (Iserenringteners) = R, T (PremiumService) = P and  $\Gamma$ (BasicService) = B. Then the Wiles of Figure 3 return he first parameter carries an With ribute of the user (its name This parameter carries an attribute of the user (its name) by a use whether it is a premium user of not. Instead, the necessary whether wish premium user or not, instead the necessary a service is premium user that checks the username and returns



### Cost-aware APP // name: lambda1.miniSL

```
// tag: premUser
( isPremiumUser, par ) => {
    if( isPremiumUser ) {
        call PremiumService( par )
    } else {
        call BasicService( par )
    }
}
```







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De Palma, G., Giallorenzo, S., Mauro, J., Trentin, M., & Zavattaro, G. (2024). FunLess: Functions-as-a-Service for Private Edge Cloud Systems. IEEE International Conference on Web Services, ICWS 2024, Shenzhen, China, July 7-13, 2024, 961–967. IEEE.







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De Palma, G., Giallorenzo, S., Mauro, J., Trentin, M., & Zavattaro, G. (2024). WebAssembly at the Edge: Benchmarking a Serverless Platform for Private Edge Cloud Systems. IEEE Internet Computing, (01), 1–8.

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#### Distributed Serverless Function Scheduling in Ad-Hoc Drone Networks Hierarchical Scheduling









#### Distributed Serverless Function Scheduling in Ad-Hoc Drone Networks Progressive Scheduling









#### Distributed Serverless Function Scheduling in Ad-Hoc Drone Networks Offloading









### **Distributed Serverless Function Scheduling in Ad-Hoc Drone Networks**



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#### **Distributed Serverless Function Scheduling in Ad-Hoc Drone Networks** AHAPP

Listing 1: ScanFireAlert policy. targets: all attributes: device: { camera, gpu }

Listing 2: AnalysePeopleScan policy. targets: all affinity: { ScanPeople } offload: ScanPeopleData



Listing 3: ScanPeople policy. targets: all attributes: position: { range: 100, lat: X, lon: Y, alt: Z } device: { camera gpu rotors: { lock: true }

*K LEONARDO* 







#### **Distributed Serverless Function Scheduling in Ad-Hoc Drone Networks** AHAPP

targets: one attributes: position: { range: 25, device: { gpu: { lock: true } } strategy: [ energy, position ]



```
Listing 4: TrackSurvivor policy.
lat: X, lon: Y, alt: Z }
camera: { lock: true }
rotors: { lock: true }
```





### Affinity-aware APP



De Palma, G., Giallorenzo, S., Mauro, J., Trentin, M., & Zavattaro, G. (2025). Affinity-aware Serverless Function Scheduling. 22nd IEEE International Conference on Software Architecture, ICSA 2025, Odense, Denmark, March 31-April 4, 2025. IEEE.



### Affinity-aware APP



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22nd IEEE International Conference on Software Architecture, ICSA 2025, Odense, Denmark, March 31-April 4, 2025. IEEE.



Deciding the reachability/ co-occurence problem	APP
Lower bound	Linear
Upper bound	Linear



Deciding the reachability/ co-occurence problem	APP	pc
Lower bound	Linear	
Upper bound	Linear	





Deciding the reachability/ co-occurence problem	APP	neg. polarised aAPP	pos. polarised aAPP
Lower bound	Linear	Linear	NP (reduction from 3SAT)
Upper bound	Linear	Linear	PSPACE (from aAPP's complexity)



Deciding the reachability/ co-occurence problem	APP	neg. polarised aAPP	pos. polarised aAPP	aAPP
Lower bound	Linear	Linear	NP (reduction from 3SAT)	PSPACE (reduction from PLANSAT)
Upper bound	Linear	Linear	PSPACE (from aAPP's complexity)	PSPACE (reduction to PLANSAT)











# A Gentle Introduction to Serverless





### Monolith



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#### Microservices









Runtime Environment





[[redirects]]
from = "/cgi-bin/hello/:name"
to = "/.netlify/functions/hello?name=:name"
status = 200





### Formalising APP's Semantics

$$[C_{start}] \xrightarrow{reg(f) = (n,t)} p(t) = \overline{b} \quad f, C, \overline{b} \to w \quad w \neq \bot \quad C(w) = (\sigma, n')$$

$$C \xrightarrow{(start,f,w)} C[w \mapsto (\sigma \cup \{f\}, n' + n, m)]$$

$$[C_{fail}] \xrightarrow{reg(f) = (\cdot,t)} p(t) = \overline{b} \quad f, C, \overline{b} \to \bot$$

$$C \xrightarrow{(fail,f)} C$$

$$[C_{done}] \xrightarrow{w \in \operatorname{dom}(C)} f \in \sigma \quad reg(f) = (n, \cdot) \quad C(w) = (\sigma, n', m)$$

$$C \xrightarrow{(done,f,w)} C[w \mapsto (\sigma \setminus \{f\}, n' - n, m)]$$

De Palma, G., Giallorenzo, S., Mauro, J., Trentin, M., & Zavattaro, G. (2024). Function-as-a-Service Allocation Policies Made Formal. In T. Margaria & B. Steffen (Eds.), Leveraging Applications of Formal Methods, Verification and Validation. REoCAS Colloquium in Honor of Rocco De Nicola - 12th International Symposium, ISoLA 2024, Crete, Greece, October 27-31, 2024, Proceedings, Part I (pp. 306–321). Springer.

$$w \neq \bot$$
  $C(w) = (\sigma, n', m)$ 

$$f, C, \overline{b} \to \bot$$





## **Formalising APP's Semantics**

De Palma, G., Giallorenzo, S., Mauro, J., Trentin, M., & Zavattaro, G. (2024). Function-as-a-Service Allocation Policies Made Formal. In T. Margaria & B. Steffen (Eds.), Leveraging Applications of Formal Methods, Verification and Validation. REoCAS Colloquium in Honor of Rocco De Nicola - 12th International Symposium, ISoLA 2024, Crete, Greece, October 27-31, 2024, Proceedings, Part I (pp. 306–321). Springer.

		Entities
	$f\in \mathcal{F}$	$w \in \mathcal{W} \subset Ident$
	$t \in \mathcal{T} \subset Identifiers$	$C \in \mathcal{C} \triangleq \mathcal{W} \to Multiset(\mathcal{F}) \times \mathbb{N}$
	$reg \in \mathcal{F}  ightarrow \mathbb{N}  imes \mathcal{T}$	$p \in \mathcal{P} \triangleq \mathcal{T} \to Li.$
$\overline{\mathrm{om}(C)}, s, \overline{i}, C \triangleright w$	$\llbracket \cdot \rrbracket : app \to \mathcal{P}$	$b \in \mathcal{B} \triangleq (List(\mathcal{W}) \cup \star) \times s_{-}opt \times List(i)$
$, C, (\star, s, \overline{i}) \to w$		
$C, b_2 :: \cdots :: b_n \to v$	$\mathcal{U}$	
$:: \cdots :: b_n \to w$		
$C(w) = (\sigma, n', m)$		
m		

m)]

 $\sigma(\sigma) = (\sigma, n', m)$ 

m)]





**Formalising APP's Semantics** Workers Layer  $[W_{first}] - \frac{\operatorname{strategy}(\overline{w}, s) = w \quad w \neq \bot \quad \operatorname{valid}(f, w, \overline{i}, C)}{f, \overline{w}, s, \overline{i}, C \triangleright w} \qquad [W_{end}] - \frac{\operatorname{strategy}(\overline{w}, s) = w}{f, \overline{w}, s, \overline{i}, C \triangleright w}$  $[W_{next}] - \frac{\mathsf{strategy}(\overline{w}, s) = w \quad w \neq \bot \quad \neg \mathsf{valid}(f, w, \overline{i}, C) \quad f, \overline{w} \setminus w, s, \overline{i}, \overline{w}, s, \overline{i}, C \triangleright w'}{f, \overline{w}, s, \overline{i}, C \triangleright w'}$ Blocks Layer  $[B_{one}] \frac{\overline{w'} = \overline{w} \cap \operatorname{dom}(C) \quad f, \overline{w'}, s, \overline{i}, C \triangleright w}{f, C, (\overline{w}, s, \overline{i}) \to w} \qquad [B_{star}] \frac{f, \overline{\operatorname{dom}(C)}, s, \overline{i}, C \triangleright w}{f, C, (\star, s, \overline{i}) \to w}$  $\begin{bmatrix} C_{start} \end{bmatrix} \xrightarrow{reg(f) = (n,t)} p(t) = b \quad f, C, b \to w \quad w \neq \bot \quad C(w) = 0$  $C \xrightarrow{(start,f,w)} C[w \mapsto (\sigma \cup \{f\}, n'+n,m)]$  $[C_{fail}] \xrightarrow{reg(f) = (\cdot, t)} p(t) = b \quad f, C, b \to \bot$  $\begin{bmatrix} C_{done} \end{bmatrix} \xrightarrow{w \in \text{dom}(C) \quad f \in \sigma \quad reg(f) = (n, \cdot) \quad C(w) = (\sigma, n', \sigma) \\ C \xrightarrow{(done, f, w)} C[w \mapsto (\sigma \setminus \{f\}, n' - n, m)] \end{bmatrix}$ 

De Palma, G., Giallorenzo, S., Mauro, J., Trentin, M., & Zavattaro, G. (2024). Function-as-a-Service Allocation Policies Made Formal. In T. Margaria & B. Steffen (Eds.), Leveraging Applications of Formal Methods, Verification and Validation. REoCAS Colloquium in Honor of Rocco De Nicola - 12th International Symposium, ISoLA 2024, Crete, Greece, October 27-31, 2024, Proceedings, Part I (pp. 306–321). Springer.

$$\mathsf{trategy}(\overline{w}, s) = \bot$$

$$f, \overline{w}, s, \overline{i}, C \triangleright \bot$$

$$s, \overline{i}, C \triangleright w'$$

	Entities
$f\in \mathcal{F}$	$w \in \mathcal{W} \subset Identi$
$t \in \mathcal{T} \subset Identifiers$	$C \in \mathcal{C} \triangleq \mathcal{W} \to Multiset(\mathcal{F}) \times \mathbb{N}$
$reg \in \mathcal{F} \to \mathbb{N}  imes \mathcal{T}$	$p \in \mathcal{P} \triangleq \mathcal{T} \to Lis$
$\llbracket \cdot \rrbracket : app \to \mathcal{P}$	$b \in \mathcal{B} \triangleq (List(\mathcal{W}) \cup \star) \times s_opt \times List(i)$









### **Cost-aware APP**



- mapReduce:
- workers:
 - wrk: W1
 - wrk: W2
 invalidate: max\_latency







#### Distributed Serverless Function Scheduling in Ad-Hoc Drone Networks Software Components • Base Station





#### Distributed Serverless Function Scheduling in Ad-Hoc Drone Networks Software Components • Node







### **Topology-aware APP** Gateway $LocalCtl_1$ $u_1$ $W_{i+1}$ #local $\overline{W_i}$ #edge $u_n$ fleet #edge (!)

De Palma, G., Giallorenzo, S., Mauro, J., Trentin, M., & Zavattaro, G. (2022). A Declarative Approach to Topology-Aware Serverless Function-Execution Scheduling. In C. A. Ardagna, ... J. Zhang (Eds.), IEEE International Conference on Web Services, ICWS 2022, Barcelona, Spain, July 10-16, 2022 (pp. 337–342). IEEE.





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## **Topology-aware APP**

critical:
<pre>- controller: LocalCtl_1</pre>
workers:
- *edge
strategy: random
followup: fail
<pre>machine_learning:</pre>
- controller: CloudCtl
workers:
<pre>- *cloud</pre>
<pre>topology_tolerance: sam</pre>
followup: default

De Palma, G., Giallorenzo, S., Mauro, J., Trentin, M., & Zavattaro, G. (2022). A Declarative Approach to Topology-Aware Serverless Function-Execution Scheduling. In C. A. Ardagna, ... J. Zhang (Eds.), IEEE International Conference on Web Services, ICWS 2022, Barcelona, Spain, July 10-16, 2022 (pp. 337–342). IEEE.

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13	default:
14	<pre>- controller: LocalCtl_</pre>
15	workers:
16	- *internal
17	strategy: random
18	- *cloud
19	strategy: random
20	<pre>strategy: best_first</pre>
21	<pre>- controller: LocalCtl_</pre>
22	workers: # same as ab
23	<pre>strategy: best_first</pre>
24	strategy: random



